32.4 Popping of Stack Frames

Recall the translation scheme for clauses:

$$
\begin{array}{lcl}
\text{code}_C\ r & = & \text{pushenv m} \\
 & & \text{code}_G\ g_1\ \rho \\
 & & \ldots \\
 & & \text{code}_G\ g_n\ \rho \\
 & & \text{popenv}
\end{array}
$$

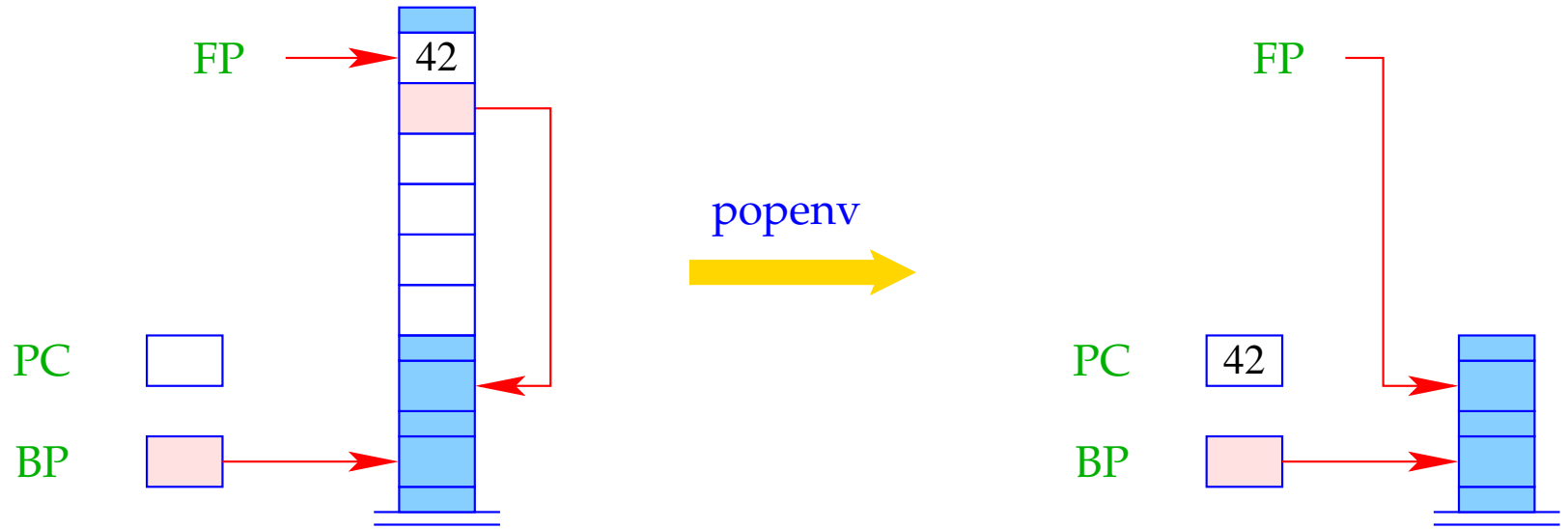
The present stack frame can be popped ...

- if the applied clause was the last (or only); and
- if all goals in the body are definitely finished.

$$\Longrightarrow \quad \text{the backtrack point is older} \quad \text{:-)}$$

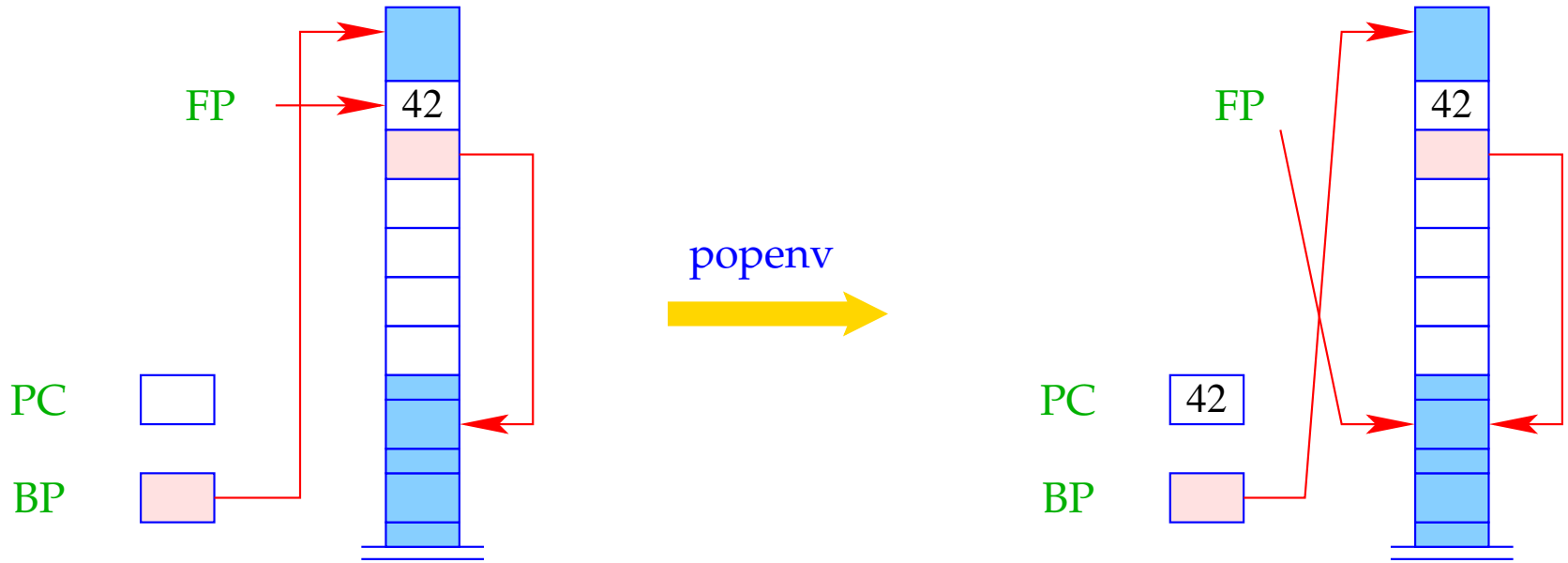
$$\Longrightarrow \quad \text{FP} > \text{BP}$$

The instruction popenv restores the registers FP and PC and possibly pops the stack frame:




```
if (FP > BP) SP = FP - 6;
PC = posCont;
FP = FPold;
```

Warning: popenv may fail to de-allocate the frame !!!




```
if (FP > BP) SP = FP - 6;
PC = posCont;
FP = FPold;
```

If popping the stack frame fails, new data are allocated on top of the stack. When returning to the frame, the locals still can be accessed through the FP :-))

33 Queries and Programs

The translation of a program: $p \equiv rr_1 \ldots rr_h ? g$
consists of:

- an instruction no for failure;
- code for evaluating the query g;
- code for the predicate definitions rr_i.

Preceding query evaluation:

$\Longrightarrow$ initialization of registers
$\Longrightarrow$ allocation of space for the globals

Succeeding query evaluation:

$\Longrightarrow$ returning the values of globals

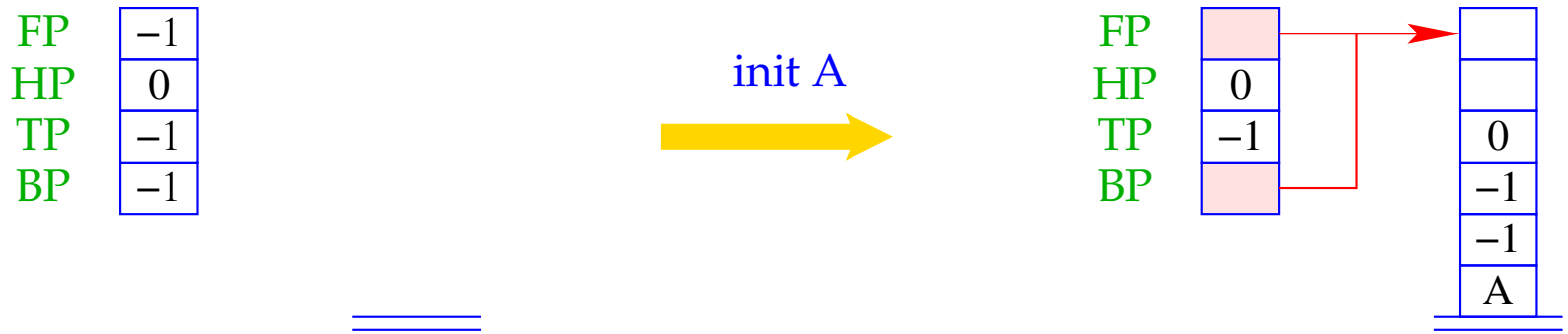
$$
\begin{array}{lll}
\text{code } p & = & \text{init A} \\
 & & \text{pushenv d} \\
 & & \text{code}_G\ g\ \rho \\
 & & \text{halt d} \\
 & \text{A:} & \text{no} \\
 & & \text{code}_P\ rr_1 \\
 & & \ldots \\
 & & \text{code}_P\ rr_h
\end{array}
$$

where $free(g) = \{X_1, \ldots, X_d\}$ and ρ is given by $\rho\ X_i = i$.

The instruction halt d ...

- ... terminates the program execution;
- ... returns the bindings of the d globals;
- ... causes backtracking — if demanded by the user :-)

The instruction init A is defined by:




```
BP = FP = SP = 5;
S[0] = A;
S[1] = S[2] = -1;
S[3] = 0;
BP = FP;
```

At address “A” for a failing goal we have placed the instruction no for printing no to the standard output and halt :-)

The Final Example:

$$t(X) \leftarrow \bar{X} = b \qquad q(X) \leftarrow s(\bar{X}) \qquad s(X) \leftarrow \bar{X} = a$$
$$p \leftarrow q(X), t(\bar{X}) \qquad s(X) \leftarrow t(\bar{X}) \qquad ? \quad p$$

The translation yields:

```
      init N          popenv          q/1:  pushenv 1       E:  pushenv 1
      pushenv 0  p/0: pushenv 1             mark D              mark G
      mark A          makr B              putref 1            putref 1
      call p/0        putvar 1            call s/1            call t/1
A:    halt 0          call q/1      D:    popenv        G:    popenv
N:    no         B:   mark C        s/1:  setbtp        F:    pushenv 1
t/1:  pushenv 1       putref 1            try E               putref 1
      putref 1        call t/1            delbtp              uatom a
      uatom b    C:   popenv              jump F              popenv
```

34 Last Call Optimization

Consider the app predicate from the beginnning:

$$\begin{aligned} \mathsf{app}(X, Y, Z) &\leftarrow X = [\,],\ Y = Z \\ \mathsf{app}(X, Y, Z) &\leftarrow X = [H|X'],\ Z = [H|Z'],\ \mathsf{app}(X', Y, Z') \end{aligned}$$

We observe:

- The recursive call occurs in the last goal of the clause.
- Such a goal is called last call.

  $\Longrightarrow$ we try to evaluate it in the current stack frame !!!

  $\Longrightarrow$ after (successful) completion, we will not return to the current caller !!!

Consider a clause r: $p(X_1, \ldots, X_k) \leftarrow g_1, \ldots, g_n$

with m locals where $g_n \equiv q(t_1, \ldots, t_h)$. The interplay between code_C and code_G:

$$
\begin{array}{rcl}
\text{code}_C\ r & = & \text{pushenv m} \\
 & & \text{code}_G\ g_1\ \rho \\
 & & \ldots \\
 & & \text{code}_G\ g_{n-1}\ \rho \\
 & & \text{mark B} \\
 & & \text{code}_A\ t_1\ \rho \\
 & & \ldots \\
 & & \text{code}_A\ t_h\ \rho \\
 & & \text{call q/h} \\
 & \text{B}: & \text{popenv}
\end{array}
$$

Replacement:			
	mark B	$\Longrightarrow$	lastmark
	call q/h; popenv	$\Longrightarrow$	lastcall q/h m

Consider a clause r: $p(X_1, \ldots, X_k) \leftarrow g_1, \ldots, g_n$

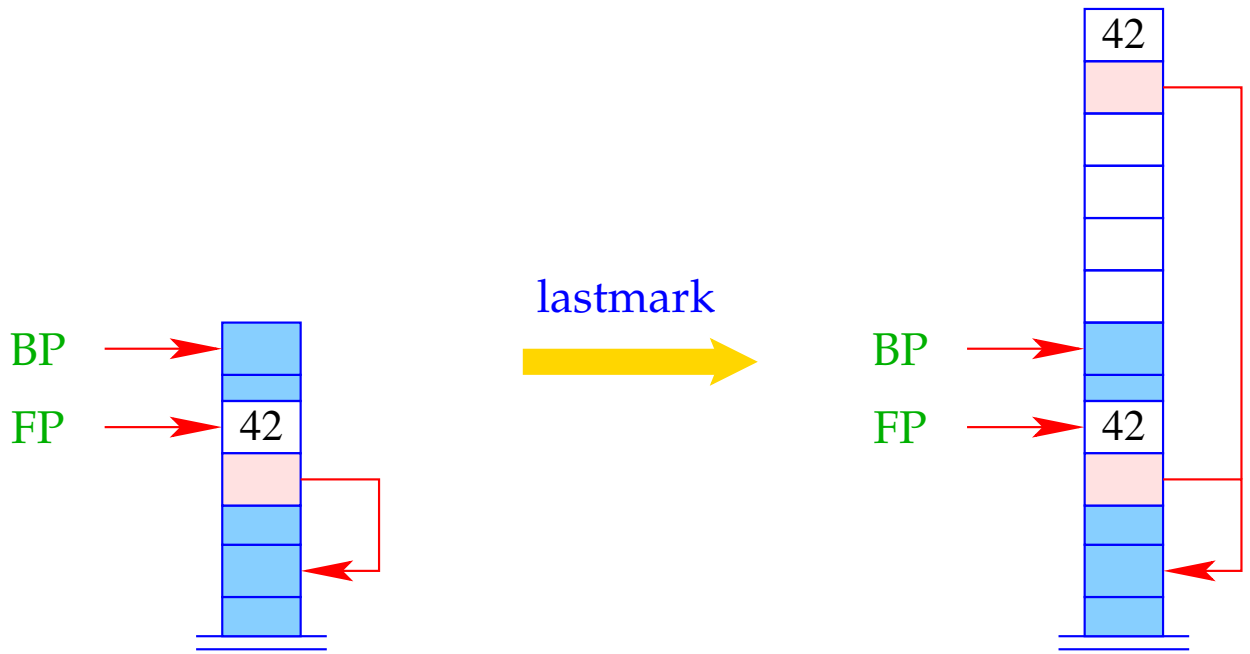
with m locals where $g_n \equiv q(t_1, \ldots, t_h)$. The interplay between code_C and code_G:

$$
\begin{array}{lll}
\text{code}_C\ r & = & \text{pushenv m} \\
 & & \text{code}_G\ g_1\ \rho \\
 & & \ldots \\
 & & \text{code}_G\ g_{n-1}\ \rho \\
 & & \text{lastmark} \\
 & & \text{code}_A\ t_1\ \rho \\
 & & \ldots \\
 & & \text{code}_A\ t_h\ \rho \\
 & & \text{lastcall q/h m}
\end{array}
$$

Replacement:			
	mark B	$\Longrightarrow$	lastmark
	call q/h; popenv	$\Longrightarrow$	lastcall q/h m

If the current clause is not last or the $g_1, \ldots, g_{n-1}$ have created backtrack points, then FP $\leq$ BP :-)

Then lastmark creates a new frame but stores a reference to the predecessor:




```
if (FP ≤ BP) {
    SP = SP + 6;
    S[SP] = posCont; S[SP-1] = FPold;
}
```

If FP $>$ BP then lastmark does nothing :-)

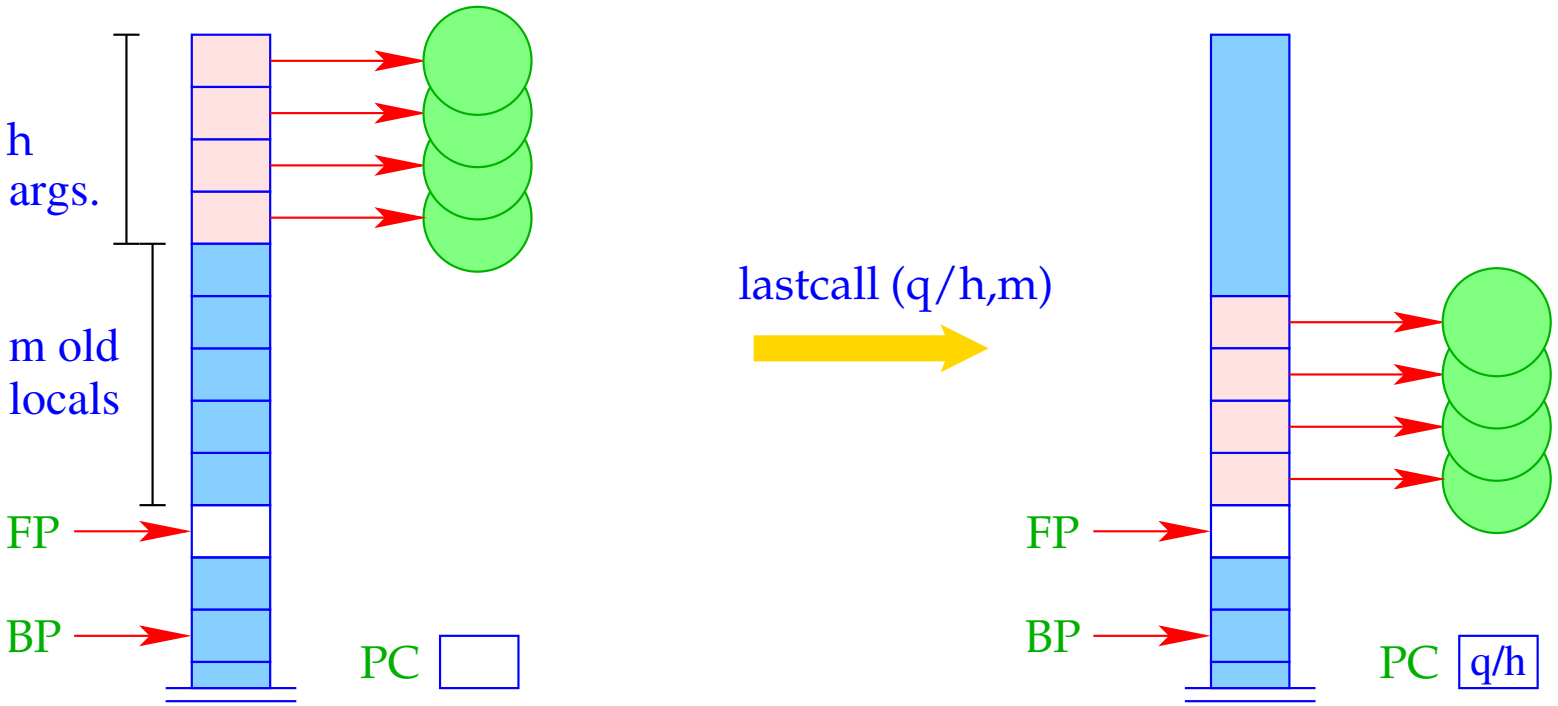
If FP $\leq$ BP, then lastcall q/h m behaves like a normal call q/h.

Otherwise, the current stack frame is re-used. This means that:

- the cells S[FP+1], S[FP+2], . . ., S[FP+h] receive the new values and
- q/h can be jumped to :-)

```
lastcall q/h m   =   if (FP ≤ BP) call q/h;
                     else {
                        move m h;
                        jump q/h;
                     }
```

The difference between the old and the new addresses of the parameters m just equals the number of the local variables of the current clause :-))


h
args.
m old
locals
FP
BP
PC
lastcall (q/h,m)
FP
BP
PC
q/h

Example:

Consider the clause:

$$\mathsf{a}(X, Y) \leftarrow \mathsf{f}(\bar{X}, X_1), \mathsf{a}(\bar{X}_1, \bar{Y})$$

The last-call optimization for $\text{code}_C\ r$ yields:

```
                    mark A          A:  lastmark
pushenv 3           putref 1            putref 3
                    putvar 3            putref 2
                    call f/2            lastcall a/2 3
```

Example:

Consider the clause:

$$\mathsf{a}(X, Y) \leftarrow \mathsf{f}(\bar{X}, X_1), \mathsf{a}(\bar{X}_1, \bar{Y})$$

The last-call optimization for $\text{code}_C\ r$ yields:

```
                    mark A          A:  lastmark
    pushenv 3       putref 1            putref 3
                    putvar 3            putref 2
                    call f/2            lastcall a/2 3
```

Note:

If the clause is last and the last literal is the only one, we can skip lastmark and can replace lastcall q/h m with the sequence move m n; jump p/n :-))

Example:

Consider the last clause of the app predicate:

$$\mathsf{app}(X, Y, Z) \quad \leftarrow \quad \bar{X} = [H|X'],\ \bar{Z} = [\bar{H}|Z'],\ \mathsf{app}(\bar{X}', \bar{Y}, \bar{Z}')$$

Here, the last call is the only one :-) Consequently, we obtain:

```
A:  pushenv 6
    putref 1
    ustruct [|]/2 B
    son 1
    uvar 4
    son 2
    uvar 5
    up C

B:  putvar 4
    putvar 5
    putstruct [|]/2
    bind

C:  putref 3
    ustruct [|]/2 D
    son 1
    uref 4
    son 2
    uvar 6
    up E

D:  check 4
    putref 4
    putvar 6
    putstruct [|]/2
    bind

E:  putref 5
    putref 2
    putref 6
    move 6 3
    jump app/3
```

35 Trimming of Stack Frames

Idea:

- Order local variables according to their life times;
- Pop the dead variables — if possible :-}

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Example:

Consider the clause:

$$\mathsf{a}(X, Z) \leftarrow \mathsf{p}_1(\bar{X}, X_1), \mathsf{p}_2(\bar{X}_1, X_2), \mathsf{p}_3(\bar{X}_2, X_3), \mathsf{p}_4(\bar{X}_3, \bar{Z})$$

35 Trimming of Stack Frames

Idea:

- Order local variables according to their life times;
- Pop the dead variables — if possible :-}

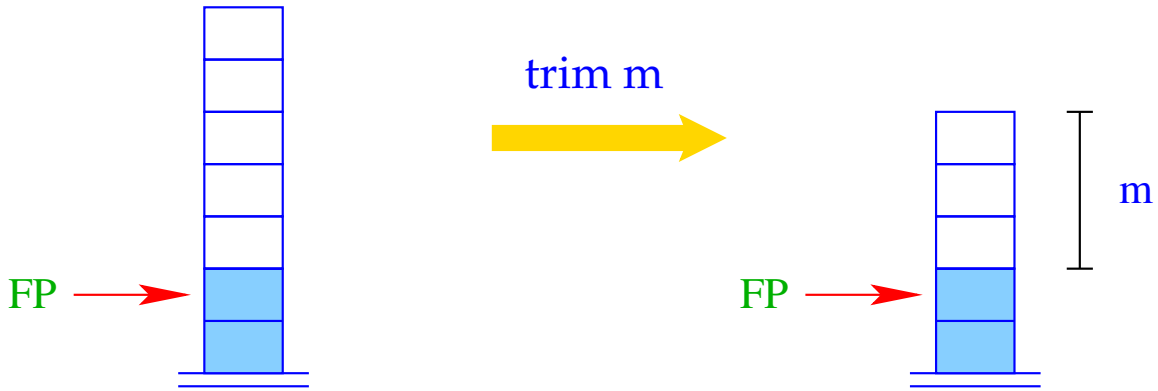
Example:

Consider the clause:

$$\mathsf{a}(X, Z) \leftarrow \mathsf{p}_1(\bar{X}, X_1), \mathsf{p}_2(\bar{X}_1, X_2), \mathsf{p}_3(\bar{X}_2, X_3), \mathsf{p}_4(\bar{X}_3, \bar{Z})$$

After the query $\mathsf{p}_2(\bar{X}_1, X_2)$, variable X_1 is dead.
After the query $\mathsf{p}_3(\bar{X}_2, X_3)$, variable X_2 is dead :-)

After every non-last goal with dead variables, we insert the instruction trim :



if (FP $\geq$ BP)
    SP = FP + m;

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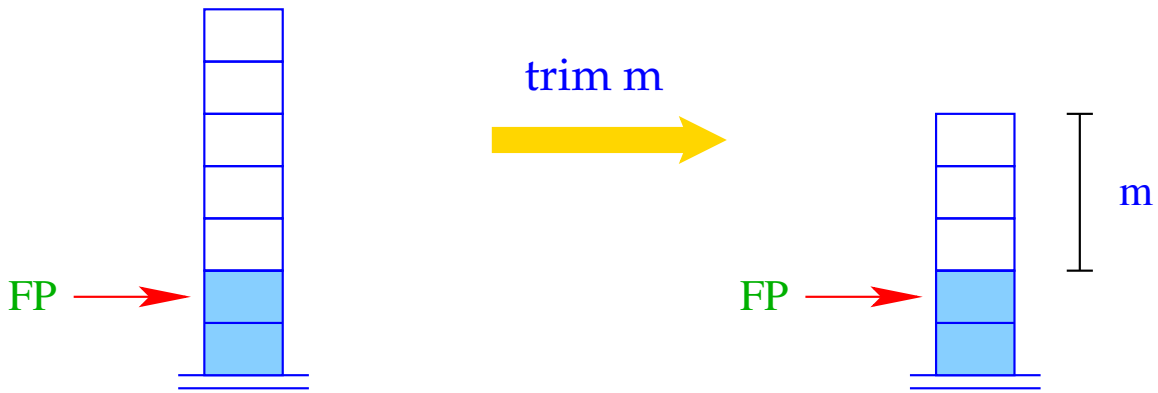


if (FP $\geq$ BP)
    SP = FP + m;

The dead locals can only be popped if no new backtrack point has been allocated :-)

Example (continued):

$$\mathsf{a}(X, Z) \leftarrow \mathsf{p}_1(\bar{X}, X_1), \mathsf{p}_2(\bar{X}_1, X_2), \mathsf{p}_3(\bar{X}_2, X_3), \mathsf{p}_4(\bar{X}_3, \bar{Z})$$

Ordering of the variables:

$$\rho = \{X \mapsto 1, Z \mapsto 2, X_3 \mapsto 3, X_2 \mapsto 4, X_1 \mapsto 5\}$$

The resulting code:

```
pushenv 5      A:  mark B          mark C      lastmark
mark A             putref 5        putref 4    putref 3
putref 1           putvar 4        putvar 3    putref 2
putvar 5           call p2/2       call p3/2   lastcall p4/2 3
call p1/2      B:  trim 4      C:  trim 3
```

36 Clause Indexing

Observation:

Often, predicates are implemented by case distinction on the first argument.

$\Longrightarrow$ Inspecting the first argument, many alternatives can be excluded :-)

$\Longrightarrow$ Failure is earlier detected :-)

$\Longrightarrow$ Backtrack points are earlier removed. :-))

$\Longrightarrow$ Stack frames are earlier popped :-)))

Example: The app-predicate:

$$\begin{aligned} \mathsf{app}(X, Y, Z) &\leftarrow X = [\,],\ Y = Z \\ \mathsf{app}(X, Y, Z) &\leftarrow X = [H|X'],\ Z = [H|Z'],\ \mathsf{app}(X', Y, Z') \end{aligned}$$

- If the root constructor is $[\,]$, only the first clause is applicable.
- If the root constructor is $[|]$, only the second clause is applicable.
- Every other root constructor should fail !!
- Only if the first argument equals an unbound variable, both alternatives must be tried ;-)

Idea:

- Introduce separate try chains for every possible constructor.
- Inspect the root node of the first argument.
- Depending on the result, perform an indexed jump to the appropriate try chain.

Assume that the predicate p/k is defined by the sequence rr of clauses $r_1 \ldots r_m$.

Let tchains rr denote the sequence of try chains as built up for the root constructors occurring in unifications $X_1 = t$.

Example:

Consider again the app-predicate, and assume that the code for the two clauses start at addresses A_1 and A_2, respectively.

Then we obtain the following four try chains:

```
VAR:   setbtp      // variables     NIL:    jump A1     // atom [ ]
       try A1
       delbtp                     CONS:   jump A2     // constructor [|]
       jump A2
                                  ELSE:   fail        // default
```

Example:

Consider again the app-predicate, and assume that the code for the two clauses start at addresses A_1 and A_2, respectively.

Then we obtain the following four try chains:

```
VAR:   setbtp      // variables     NIL:    jump A1     // atom [ ]
       try A1
       delbtp                     CONS:   jump A2     // constructor [|]
       jump A2
                                  ELSE:   fail        // default
```

The new instruction fail takes care of any constructor besides [ ] and [|] ...

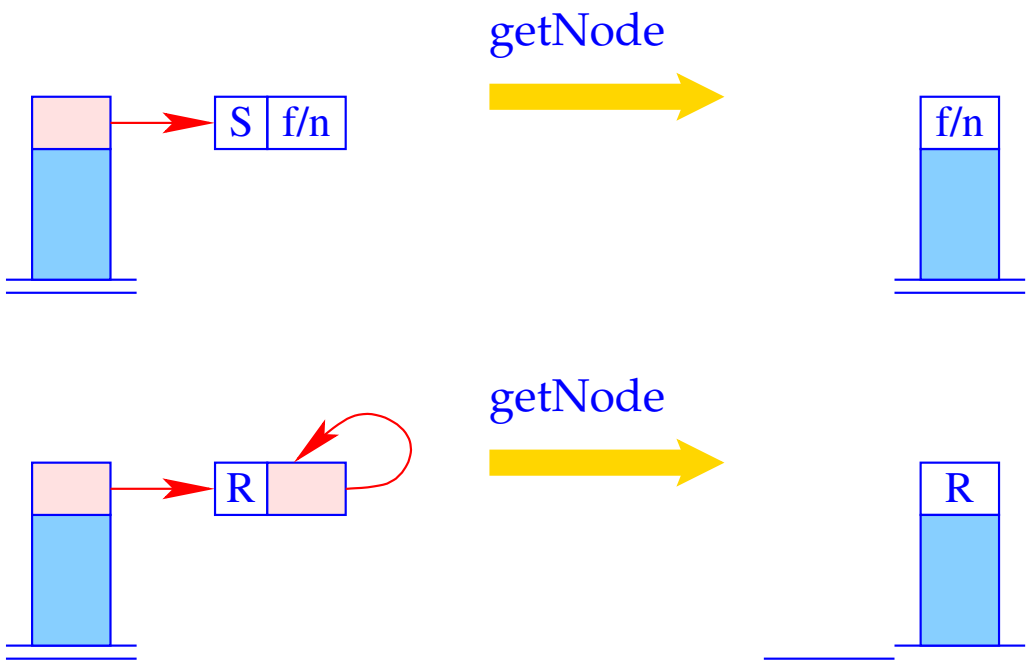
```
fail  =  backtrack()
```

It directly triggers backtracking :-)

Then we generate for a predicate p/k:

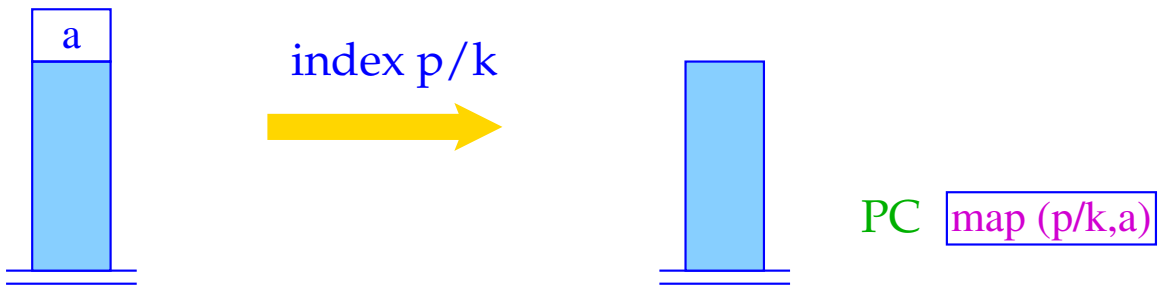
$$
\begin{array}{lll}
\text{code}_P\ rr = & \text{putref 1} & \\
 & \text{getNode} & \text{// extracts the root label} \\
 & \text{index p/k} & \text{// jumps to the try block} \\
 & \text{tchains } rr & \\
A_1: & \text{code}_C\ r_1 & \\
 & \ldots & \\
A_m: & \text{code}_C\ r_m &
\end{array}
$$

The instruction getNode returns "R" if the pointer on top of the stack points to an unbound variable. Otherwise, it returns the content of the heap object:



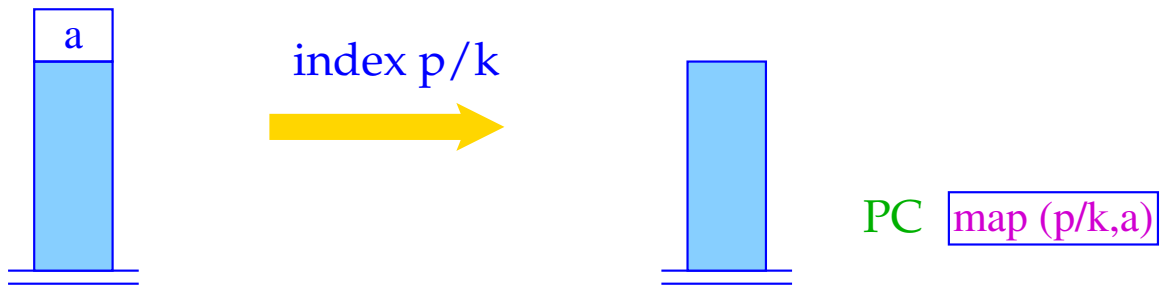

```
switch (H[S[SP]]) {
case (S, f/n):    S[SP] = f/n; break;
case (A,a):       S[SP] = a; break;
case (R,_) :      S[SP] = R;
}
```

The instruction index p/k performs an indexed jump to the appropriate try chain:




```
PC = map (p/k,S[SP]);
SP– –;
```

The instruction index p/k performs an indexed jump to the appropriate try chain:



PC = map (p/k,S[SP]);
SP– –;

The function map() returns, for a given predicate and node content, the start address of the appropriate try chain :-)

It typically is defined through some hash table :-))