Technische Universität München Fakultät für Informatik Prof. Dr. H. Seidl Winter Semester 10/11 Vesal Vojdani Aleksandr Karbyshev

## **Program Optimization**

## Solutions to Sheet 3

Exercise 1: (P/H) Height of a lattice

10 Points

For a lattice  $\mathbb{D}$ , we define the height of the lattice  $h(\mathbb{D}) = n$  as the maximal length of strictly ascending chains  $d_0 \sqsubset d_1 \sqsubset d_2 \sqsubset \cdots \sqsubset d_n$  in the lattice. (Note that finiteness of height does not imply finiteness of a lattice. Give an example!)

Let  $\mathbb{D}_1, \mathbb{D}_2, \mathbb{D}$  be lattices of finite height. Show that the following hold:

- a) (P)  $h(\mathbb{D}_1 \times \mathbb{D}_2) = h(\mathbb{D}_1) + h(\mathbb{D}_2);$
- b) (H)  $h(\mathbb{D}^k) = k \cdot h(\mathbb{D})$ , for any  $k \in \mathbb{N}$ ;
- c) (H)  $h(X \to \mathbb{D}) = |X| \cdot h(\mathbb{D})$ , where |X| is a cardinality of the finite set X;
- d) (H)  $h([\mathbb{D}_1 \to \mathbb{D}_2]) = |\mathbb{D}_1| \cdot h(\mathbb{D}_2)$ , where  $\mathbb{D}_1$  is finite.
- a) For,  $h(\mathbb{D}_1 \times \mathbb{D}_2) = h(\mathbb{D}_1) + h(\mathbb{D}_2)$ , we abbreviate h, n and m, for the heights of  $\mathbb{D}_1 \times \mathbb{D}_2$ ,  $\mathbb{D}_1$ , and  $\mathbb{D}_2$ , respectively. We want to prove that h = n + m by showing, first,  $h \ge n + m$  and, second,  $h \le n + m$ .
  - 1. We know the sequence  $\bot \sqsubseteq a_1 \sqsubseteq \cdots \sqsubseteq a_n$  exists in  $\mathbb{D}_1$  and  $\bot \sqsubseteq b_1 \sqsubseteq \cdots \sqsubseteq b_m$  exists in  $\mathbb{D}_2$ . We form the sequence  $(\bot, \bot) \sqsubseteq (a_1, \bot) \sqsubseteq \cdots \sqsubseteq (a_n, \bot) \sqsubseteq (a_n, b_1) \sqsubseteq \cdots \sqsubseteq (a_n, b_m)$ , which shows that  $h \ge n + m$ .
  - 2. Take an ascending chain  $(a_0, b_0) \sqsubset \cdots \sqsubset (a_k, b_k)$ . Now, a simple induction shows that  $|\{a_0, \ldots, a_k\}| + |\{b_0, \ldots, b_k\}| \ge k + 2$ , because each strict inequality requires at least one element to differ. If we look at the elements  $A = \{a_0, \ldots, a_k\}$  and  $B = \{b_0, \ldots, b_k\}$ , we see that these form ascending chains in  $\mathbb{D}_1$  and  $\mathbb{D}_2$ , respectively. We obtain  $|A| \le n + 1$  and  $|B| \le m + 1$ . (We need to add one, because a sequence of length k contains k + 1 elements.) Putting it together, we have  $k + 2 \le m + n + 2$ , and  $k \le m + n$ .
- b) Use induction, Luke.
- c) You can either prove it directly, or reduce it to b) by showing that for the finite X there is an isomorphism between  $X \to \mathbb{D}$  and a product lattice.

A monotonic analysis framework is triple  $\langle \mathbb{D}, \mathcal{F}, tf \rangle$ , where  $\mathbb{D} = \langle 2^X, \sqsubseteq \rangle$  is the complete lattice of subsets of a finite set X (here,  $\sqsubseteq = \subseteq$ ),  $\mathcal{F} \subseteq [\mathbb{D} \to \mathbb{D}]$  is a set of monotonic functions, and  $tf : E \to \mathcal{F}$  is a mapping from edges in the control flow graph to transfer functions from  $\mathcal{F}$ . (This mapping tf was denoted in the lecture by  $\llbracket - \rrbracket^{\sharp}$ , i.e.,  $\llbracket k \rrbracket^{\sharp} \in \mathcal{F}$ , for any  $k \in E$ .)

For a set  $Y \subseteq X$ , we denote its complement as  $\overline{Y} = X \setminus Y$ . Now, we define the complement of a monotonic analysis framework  $\langle \mathbb{D}, \mathcal{F}, tf \rangle$  as the triple  $\langle \overline{\mathbb{D}}, \overline{\mathcal{F}}, \overline{tf} \rangle$ , where

$$\begin{array}{ll} \overline{\mathbb{D}} = \langle 2^X, \rightrightarrows \rangle & \text{the lattice ordering is } \textit{reversed}, \text{ i.e., } \sqsubseteq_{\overline{\mathbb{D}}} = \supseteq, \\ \overline{f}(Y) = \overline{f(\overline{Y})} & \text{complement of a function,} \\ \overline{\mathcal{F}} = \{\overline{f} \mid f \in \mathcal{F}\} & \text{complement of a set of functions,} \\ \overline{tf}_k = \overline{tf}_k & \text{complement of the transfer function, } k \in E. \end{array}$$

- a) Show that the complement of a monotonic analysis framework is itself a monotonic analysis framework.
- b) Let  $\mathcal{A}[u]$  (for all program points u) be a solution to the system

$$tf_k(\mathcal{A}[u]) \sqsubseteq \mathcal{A}[v] \quad (k = (u, \cdot, v) \in E).$$

Show that its complement  $\overline{A[u]}$  is a solution to the complementary analysis.

c) The complement of the Possibly Live Variables analysis described in the lecture is a Definitely Dead Variables analysis. Describe the lattice for this analysis and give the transfer function for assignments.