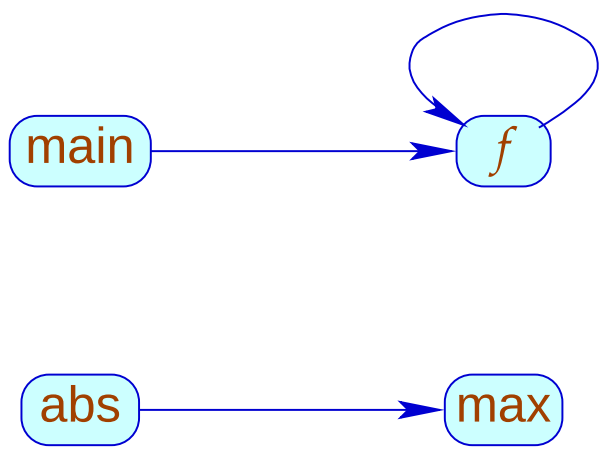
Detection of Recursion:

We construct the call-graph of the program.

In the Examples:



Call-Graph:

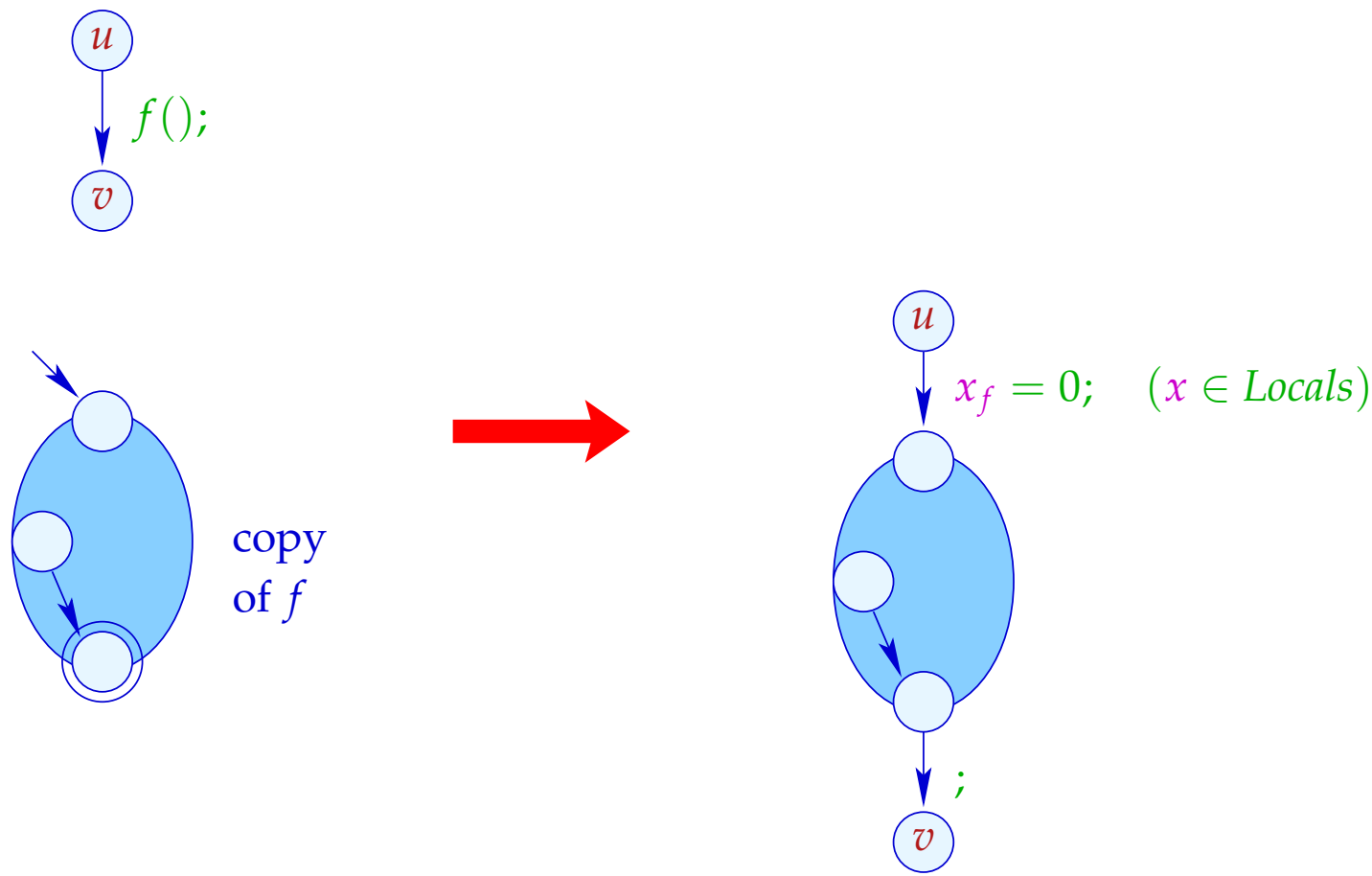
- The nodes are the procedures.
- An edge connexts g with h, whenever the body of g contains a call of h.

Strategies for Inlining:

- Just copy nur leaf-procedures, i.e., procedures without further calls :-)
- Copy all non-recursive procedures!

... here, we consider just leaf-procedures ;-)

Transformation 9:



Note:

- The Nop-edge can be eliminated if the *stop*-node of f has no out-going edges ...
- The x_f are the copies of the locals of the procedure f.
- According to our semantics of procedure calls, these must be initialized with 0 :-)

2. Idea: Elimination of Tail Recursion

```
f () { int b;
       if (a2 ≤ 1) { ret = a1; goto _exit; }
       b = a1 · a2;
       a2 = a2 − 1;
       a1 = b;
       f ();
_exit :
}
```

After the procedure call, nothing in the body remains to be done.

$\Longrightarrow$ We may directly jump to the beginning :-)

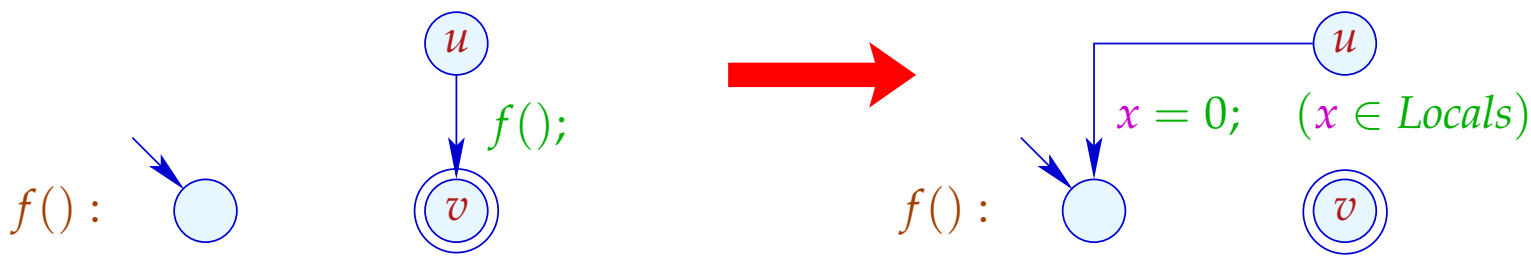
... after having reset the locals to 0.

... this yields in the Example:

$$
\begin{array}{ll}
f\,()\ \{ & \text{int}\ \ b; \\
\_f: & \text{if}\ \ (a_2 \leq 1)\ \ \{\ \ \text{ret} = a_1;\ \text{goto}\ \_exit;\ \ \} \\
 & b = a_1 \cdot a_2; \\
 & a_2 = a_2 - 1; \\
 & a_1 = b; \\
 & b = 0;\ \ \text{goto}\ \ \_f; \\
\_exit: & \\
\} &
\end{array}
$$

// It works, since we have ruled out references to variables!

Transformation 11:



Warning:

- → This optimization is crucial for programming languages without iteration constructs !!!
- → Duplication of code is not necessary :-)
- → No variable renaming is necessary :-)
- → The optimization may also be profitable for non-recursive tail calls :-)
- → The corresponding code may contain jumps from the body of one procedure into the body of another ???

Background 4: Interprocedural Analysis

So far, we can analyze each procedure separately.

→ The costs are moderate :-)

→ The methods also work in presence of separate compilation :-)

→ At procedure calls, we must assume the worst case :-(

→ Constant propagation only works for local constants :-((

Question:

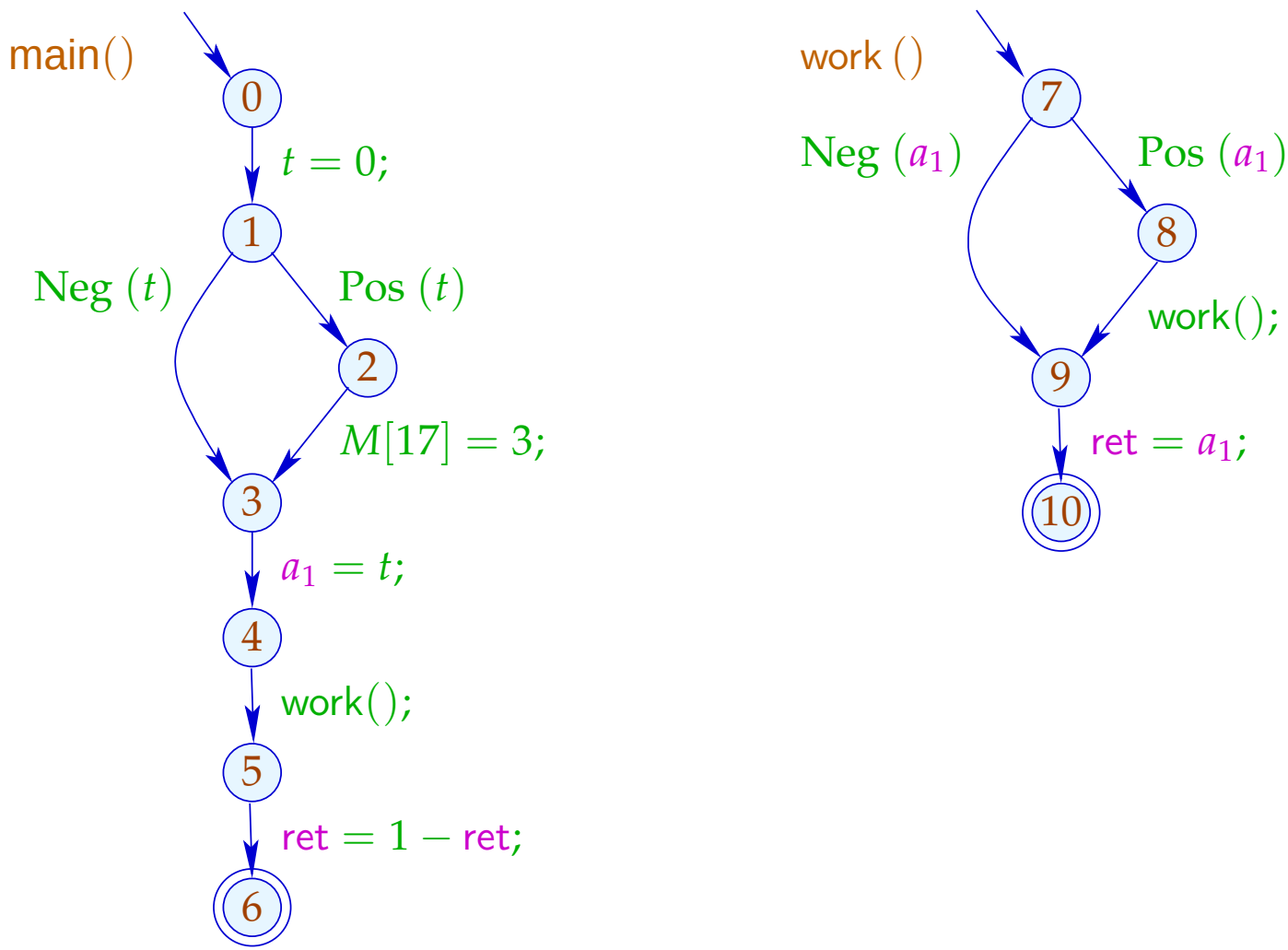
How can recursive programs be analyzed ???

Example: Constant Propagation

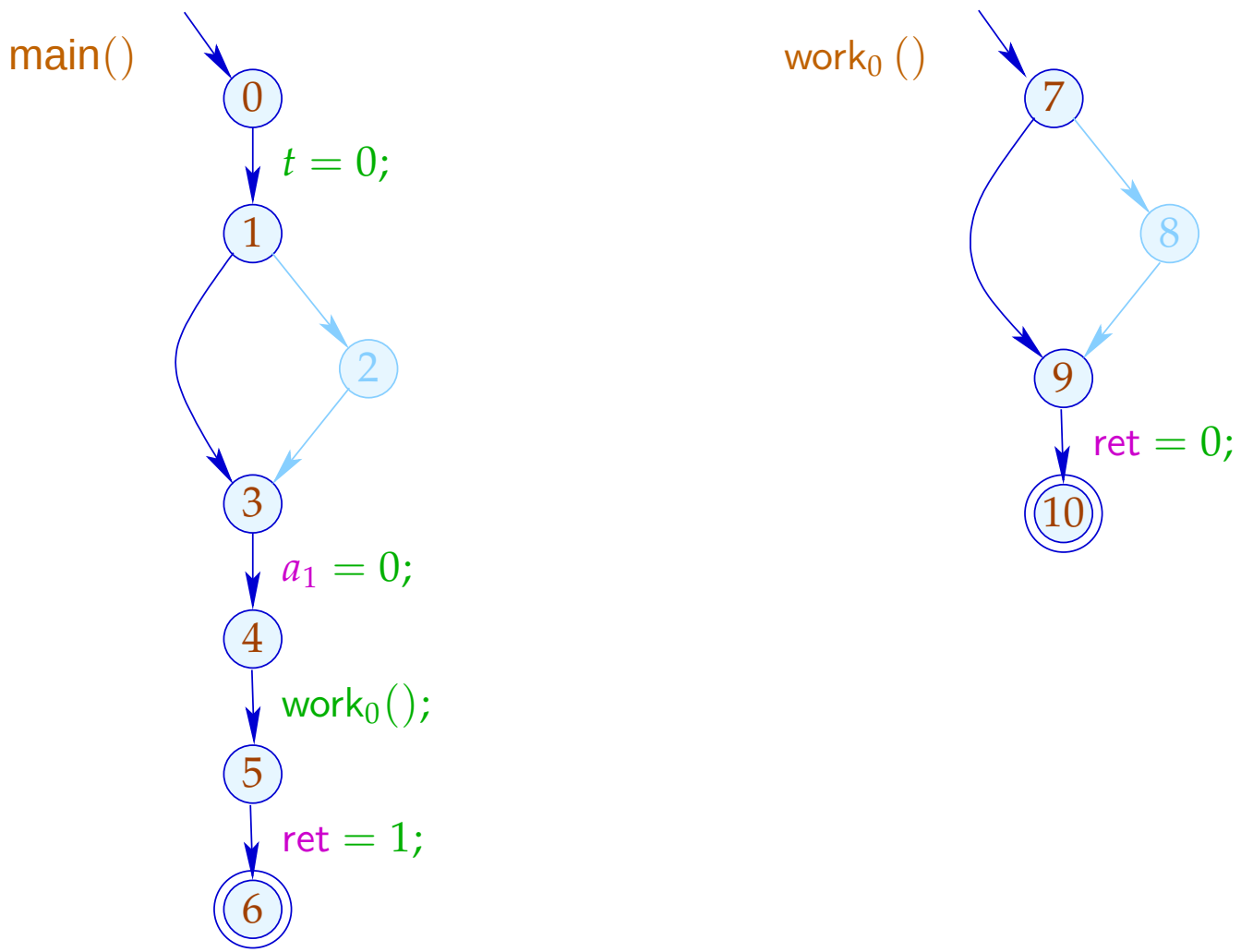
```
main() { int t;
        t = 0;
        if (t) M[17] = 3;
        a1 = t;
        work();
        ret = 1 − ret;
}
```

```
work() {
        if (a1) work();
        ret = a1;
}
```

Example: Constant Propagation



Example: Constant Propagation



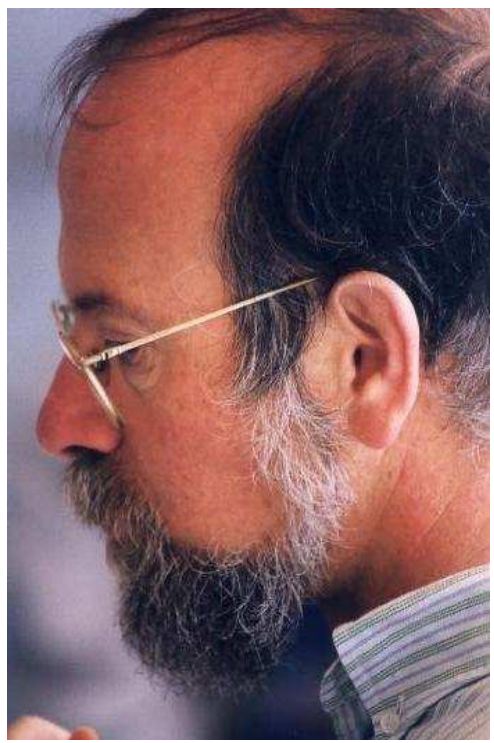
(1) Functional Approach:

Let $\mathbb{D}$ denote a complete lattice of (abstract) states.

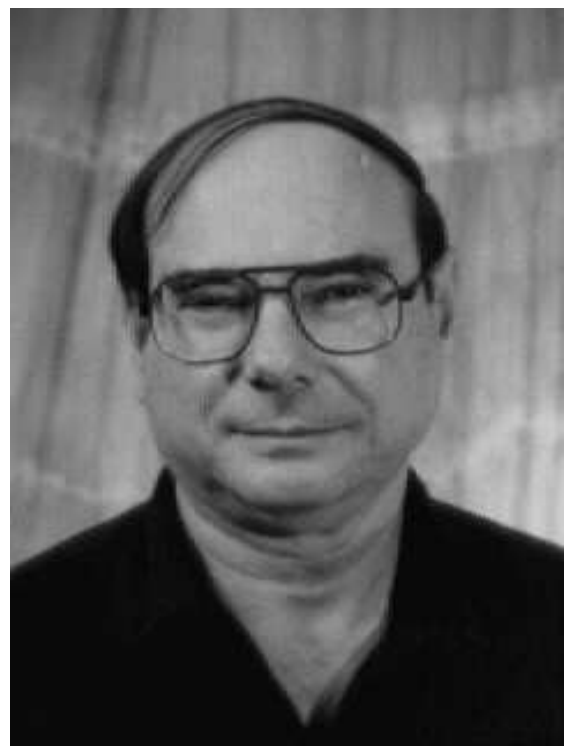
Idea:

Represent the effect of $f()$ by a function:

$$[\![f]\!]^\sharp \ : \ \mathbb{D} \to \mathbb{D}$$



Micha Sharir, Tel Aviv University



Amir Pnueli, Weizmann Institute

In order to determine the effect of a call edge $k = (u, f\ ();, v)$ we require abstract functions:

$$\begin{array}{lcl} \mathsf{enter}^\sharp & : & \mathbb{D} \to \mathbb{D} \\ \mathsf{combine}^\sharp & : & \mathbb{D}^2 \to \mathbb{D} \end{array}$$

Then we define:

$$[\![k]\!]^\sharp\ D \quad = \quad \mathsf{combine}^\sharp\ (D, [\![f]\!]^\sharp\ (\mathsf{enter}^\sharp\ D))$$

... for Constant Propagation:

$$\mathbb{D} = (\mathit{Vars} \to \mathbb{Z}^{\top})_{\bot}$$

$$\mathsf{enter}^{\sharp}\ D = \begin{cases} \bot & \text{if} \quad D = \bot \\ D|_{\mathit{Globals}} \oplus \{x \mapsto 0 \mid x \in \mathit{Locals}\} & \text{otherwise} \end{cases}$$

$$\mathsf{combine}^{\sharp}\ (D_1, D_2) = \begin{cases} \bot & \text{if} \quad D_1 = \bot \vee D_2 = \bot \\ D_1|_{\mathit{Locals}} \oplus D_2|_{\mathit{Globals}} & \text{otherwise} \end{cases}$$

The effects $[\![f]\!]^\sharp$ then can be determined by a system of constraints over the complete lattice $\mathbb{D} \to \mathbb{D}$:

$$
\begin{array}{lll}
[\![v]\!]^\sharp & \sqsupseteq & \mathsf{Id} \qquad\qquad v \quad \text{entry point} \\
[\![v]\!]^\sharp & \sqsupseteq & [\![k]\!]^\sharp \circ [\![u]\!]^\sharp \qquad k = (u, \_, v) \quad \text{edge} \\
[\![f]\!]^\sharp & \sqsupseteq & [\![stop_f]\!]^\sharp \qquad stop_f \quad \text{end point of} \quad f
\end{array}
$$

$[\![v]\!]^\sharp$: $\mathbb{D} \to \mathbb{D}$ describes the effect of all prefixes of computation forests w of a procedure which lead from the entry point to v :-)

Problems:

- How can we represent functions $f : \mathbb{D} \to \mathbb{D}$???
- If $\#\mathbb{D} = \infty$, then $\mathbb{D} \to \mathbb{D}$ has infinite strictly increasing chains :-(

Simplification: Copy-Constants

$\to$ Conditions are interpreted as ; :-)

$\to$ Only assignments $x = e;$ with $e \in \mathit{Vars} \cup \mathbb{Z}$ are treated exactly :-)

Observation:

→ The effects of assignments are:

$$[\![x = e;]\!]^\sharp\, D \quad = \quad \begin{cases} D \oplus \{x \mapsto c\} & \text{if} \quad e = c \in \mathbb{Z} \\ D \oplus \{x \mapsto (D\ y)\} & \text{if} \quad e = y \in \mathit{Vars} \\ D \oplus \{x \mapsto \top\} & \text{otherwise} \end{cases}$$

→ Let $\mathbb{V}$ denote the (finite !!!) set of constant right-hand sides. Then variables may only take values from $\mathbb{V}^\top$:-))

→ The occurring effects can be taken from

$$\mathbb{D}_f \to \mathbb{D}_f \qquad \text{with} \qquad \mathbb{D}_f = (\mathit{Vars} \to \mathbb{V}^\top)_\perp$$

→ The complete lattice is huge, but finite !!!

Improvement:

$\rightarrow$ Not all functions from $\mathbb{D}_f \rightarrow \mathbb{D}_f$ will occur :-)

$\rightarrow$ All occurring functions $\lambda D. \bot \neq M$ are of the form:

$$\begin{array}{lcll} M & = & \{x \mapsto (b_x \sqcup \bigsqcup_{y \in I_x} y) \mid x \in \mathit{Vars}\} & \text{where:} \\ M\ D & = & \{x \mapsto (b_x \sqcup \bigsqcup_{y \in I_x} D\ y) \mid x \in \mathit{Vars}\} & \text{für} \quad D \neq \bot \end{array}$$

$\rightarrow$ Let $\mathbb{M}$ denote the set of all these functions. Then for $M_1, M_2 \in \mathbb{M}$ $(M_1 \neq \lambda D. \bot \neq M_2)$:

$$(M_1 \sqcup M_2)\ x \quad = \quad (M_1\ x) \sqcup (M_2\ x)$$

$\rightarrow$ For $k = \#\mathit{Vars}$, $\mathbb{M}$ has height $\mathcal{O}(k^2)$:-)